
List of Notation

\circ	entry-wise multiplication of vectors
∇	the gradient: $\nabla f(x) = (D_1 f(x), \dots, D_n f(x))$
\neg	logical NOT
\exists	$S \ni i$ is equivalent to $i \in S$
\oplus	logical XOR (exclusive-or)
$\ \hat{f}\ _p$	$(\sum_{\gamma \in \mathbb{F}_2^n} \hat{f}(\gamma) ^p)^{1/p}$
Δ	symmetric difference of sets; i.e., $S \Delta T = \{i : i \text{ is in exactly one of } S, T\}$
\vee	logical OR
\wedge	logical AND
$*$	the convolution operator
$[z^k]F(z)$	coefficient on z^k in the power series $F(z)$
1_A	0-1 indicator function for A
$\mathbf{1}_B$	0-1 indicator random variable for event B
1_x	the 0-1 indicator function for x
2^A	the set of all subsets of A
AND_n	the logical AND function on n bits: False unless all inputs are True
$\text{Aut}(f)$	the group of automorphisms of Boolean function f
$\text{BitsToGaussians}_M^d$	on input the bit matrix $x \in \{-1, 1\}^{d \times M}$, has output $z \in \mathbb{R}^d$ equal to $\frac{1}{\sqrt{M}}$ times the column-wise sum x ; if d is omitted it's taken to be 1
\mathbb{C}	the complex numbers

$\chi(b)$	when $b \in \mathbb{F}_2^n$, denotes $(-1)^b \in \mathbb{R}$
$\chi_S(x)$	when $x \in \mathbb{R}^n$, denotes $\prod_{i \in S} x_i$, where $S \subseteq [n]$; when $x \in \mathbb{F}_2^n$, denotes $(-1)^{\sum_{i \in S} x_i}$ if $x \in \mathbb{F}_2^n$
$\text{codim} H$	for subspace $H \leq \mathbb{F}^n$, denotes $n - \dim H$
$\mathbf{Cov}[f, g]$	the covariance of f and g , $\mathbf{Cov}[f] = \mathbf{E}[fg] - \mathbf{E}[f]\mathbf{E}[g]$
$d_{\chi^2}(\varphi, 1)$	chi-squared distance of the distribution with density φ from the uniform distribution
$\text{deg}(f)$	the degree of f ; the least k such that f is a real linear combination of k -juntas
$\text{deg}_{\mathbb{F}_2}(f)$	for Boolean-valued f , the degree of its \mathbb{F}_2 -polynomial representation
$\Delta^{(\pi)}(f)$	the expected number of queries made by the best decision tree computing f when the input bits are chosen from the distribution π
$\delta^{(\pi)}(f)$	the revelation of f ; i.e., $\min\{\max_i \delta_i^{(\pi)}(\mathcal{T}) : \mathcal{T} \text{ computes } f\}$
$\Delta^{(\pi)}(\mathcal{T})$	the expected number of queries made by randomized decision tree \mathcal{T} when the input bits are chosen from the distribution π
$\delta_i^{(\pi)}(\mathcal{T})$	the probability randomized decision tree \mathcal{T} queries coordinate i when the input bits are chosen from the distribution π
$\Delta_y f$	for $f : \mathbb{F}_2^n \rightarrow \mathbb{F}_2$, the function $\mathbb{F}_2^n \rightarrow \mathbb{F}_2$ defined by $\Delta_y f(x) = f(x + y) - f(x)$
D_i	the i th discrete derivative: $D_i f(x) = \frac{f(x^{(i-1)}) - f(x^{(i--1)})}{2}$
$\text{dist}(g, h)$	the relative Hamming distance; i.e., the fraction of places where g and h disagree
$\Delta(x, y)$	the Hamming distance, $\#\{i : x_i \neq y_i\}$
$\text{DNF}_{\text{size}}(f)$	least possible size of a DNF formula computing f
$\text{DNF}_{\text{width}}(f)$	least possible width of a DNF formula computing f
$\text{DT}(f)$	least possible depth of a decision tree computing f
$\text{DT}_{\text{size}}(f)$	least possible size of a decision tree computing f
$d_{\text{TV}}(\varphi, \psi)$	total variation distance between the distributions with densities φ, ψ
\mathbf{E}_i	the i th expectation operator: $\mathbf{E}_i f(x) = \mathbf{E}_{x_i}[f(x_1, \dots, x_{i-1}, x_i, x_{i+1}, \dots, x_n)]$
\mathbf{E}_I	the expectation over coordinates I operator
$\mathbf{Ent}[f]$	for a nonnegative function on a probability space, denotes $\mathbf{E}[f \ln f] - \mathbf{E}[f] \ln \mathbf{E}[f]$

$\mathbf{E}_{\pi_p}[\cdot]$	an abbreviation for $\mathbf{E}_{\mathbf{x} \sim \pi_p^{\otimes n}}[\cdot]$
$f \oplus g$	if $f : \{-1, 1\}^m \rightarrow \{-1, 1\}$ and $g : \{-1, 1\}^n \rightarrow \{-1, 1\}$, denotes the function $h : \{-1, 1\}^{m+n} \rightarrow \{-1, 1\}$ defined by $h(x, y) = f(x)g(y)$
$f \otimes g$	if $f : \{-1, 1\}^m \rightarrow \{-1, 1\}$ and $g : \{-1, 1\}^n \rightarrow \{-1, 1\}$, denotes the function $h : \{-1, 1\}^{mn} \rightarrow \{-1, 1\}$ defined by $h(x^{(1)}, \dots, x^{(m)}) = f(g(x^{(1)}), \dots, g(x^{(m)}))$
f^{*n}	the n -fold convolution, $f * f * \dots * f$
f^\dagger	the Boolean dual defined by $f^\dagger(x) = -f(-x)$
f^{+z}	if $f : \mathbb{F}_2^n \rightarrow \mathbb{R}$, $z \in \mathbb{F}_2^n$, denotes the function $f^{+z}(x) = f(x+z)$
f_H^{+z}	denotes $(f^{+z})_H$
\mathbb{F}_2	the finite field of size 2
$\widehat{\mathbb{F}_2^n}$	the group (vector space) indexing the Fourier characters of functions $f : \mathbb{F}_2^n \rightarrow \mathbb{R}$
f^{even}	the even part of f , $(f(x) + f(-x))/2$
$\langle f, g \rangle$	$\mathbf{E}_{\mathbf{x}}[f(\mathbf{x})g(\mathbf{x})]$
f_H	if $f : \mathbb{F}_2^n \rightarrow \mathbb{R}$, $H \leq \mathbb{F}_2^n$, denotes the restriction of f to H
$\widehat{f}(i)$	shorthand for $\widehat{f}(\{i\})$ when $i \in \mathbb{N}$
$f^{\subseteq J}$	$\sum_{S \subseteq J} \widehat{f}(S) \chi_S$
$f _z$	if $f : \{-1, 1\}^n \rightarrow \mathbb{R}$, $J \subseteq [n]$, and $z \in \{-1, 1\}^{\overline{J}}$, denotes the restriction of f given by fixing the coordinates in \overline{J} to z
$f_{J z}$	if $f : \{-1, 1\}^n \rightarrow \mathbb{R}$, $J \subseteq [n]$, and $z \in \{-1, 1\}^{\overline{J}}$, denotes the restriction of f given by fixing the coordinates in \overline{J} to z
$f^{=k}$	$\sum_{ S =k} \widehat{f}(S) \chi_S$
$f^{\leq k}$	$\sum_{ S \leq k} \widehat{f}(S) \chi_S$
f^{odd}	the odd part of f , $(f(x) - f(-x))/2$
\mathbb{F}_{p^ℓ}	for p prime and $\ell \in \mathbb{N}^+$, denotes the finite field of p^ℓ elements
$\widehat{f}(S)$	the Fourier coefficient of f on character χ_S
$\mathbf{F}_{S J} \widehat{f}(z)$	for $S \subseteq J \subseteq [n]$, denotes $\widehat{f}_{J z}(S)$
$f^{\otimes d}$	if $f : \{-1, 1\}^n \rightarrow \{-1, 1\}$ then $f^{\otimes d} : \{-1, 1\}^{n^d} \rightarrow \{-1, 1\}$ is defined inductively by $f^{\otimes 1} = f$, $f^{\otimes(d+1)} = f \otimes f^{\otimes d}$
\widetilde{f}	the randomization/symmetrization of f , defined by $\widetilde{f}(r, x) = \sum_S \mathbf{r}^S f^S(x)$
$\gamma^+(\partial A)$	the Gaussian Minkowski content of ∂A

$\mathcal{G}(v, p)$	the Erdős–Rényi random graph distribution, $\pi_p^{\otimes \binom{v}{2}}$
h_i	the i th (normalized) Hermite polynomial, $h_i = \frac{1}{\sqrt{i!}} H_i$
h_α	for $\alpha \in \mathbb{N}^n$ a multi-index, the n -variate (normalized) Hermite polynomial $h_\alpha(z) = \prod_{i=1}^n h_{\alpha_i}(z_i)$
H_i	the i th probabilists' Hermite polynomial, defined by $\exp(tz - \frac{1}{2}t^2) = \sum_{i=0}^{\infty} \frac{1}{i!} H_i(z) t^i$
A^\perp	$\{\gamma : \gamma \cdot x = 0 \text{ for all } x \in A\}$
$\mathbf{Inf}_i[f]$	the influence of coordinate i on f
$\mathbf{Inf}_i^{(\rho)}[f]$	the ρ -stable influence, $\mathbf{Stab}_\rho[D_i f]$
$\widetilde{\mathbf{Inf}}_J[f]$	the coalitional influence of $J \subseteq [n]$ on $f : \{-1, 1\}^n \rightarrow \{-1, 1\}$, namely $\mathbf{Pr}_{\mathbf{z} \sim \{-1, 1\}^{\bar{J}}} [f_{J \mathbf{z}} \text{ is not constant}]$
$\widetilde{\mathbf{Inf}}_J^b[f]$	for $b \in \{-1, 1\}$, equals $\mathbf{Pr}_{\mathbf{z} \sim \{-1, 1\}^{\bar{J}}} [f_{J \mathbf{z}} \neq -b] - \mathbf{Pr}[f = b]$
\bar{J}	if $J \subseteq [n]$, denotes $[n] \setminus J$
$L^2(\{-1, 1\}^n)$	denotes $L^2(\{-1, 1\}^n, \pi_{1/2}^{\otimes n})$
$L^2(G^n)$	if G is a finite abelian group, denotes the complex inner product space of functions $G^n \rightarrow \mathbb{R}$ with inner product $\langle f, g \rangle = \mathbf{E}_{\mathbf{x} \sim G^n} [f(\mathbf{x}) \overline{g(\mathbf{x})}]$
$L^2(\Omega, \pi)$	the inner product space of (square-integrable) functions $\Omega \rightarrow \mathbb{R}$ with inner product $\langle f, g \rangle = \mathbf{E}_{\mathbf{x} \sim \pi} [f(\mathbf{x}) g(\mathbf{x})]$
$\Lambda_\rho(\alpha)$	$\mathbf{Pr}[\mathbf{z}_1 \leq t, \mathbf{z}_2 \leq t]$, where $\mathbf{z}_1, \mathbf{z}_2$ are standard Gaussians with correlation $\mathbf{E}[\mathbf{z}_1 \mathbf{z}_2] = \rho$ and $t = \Phi^{-1}(\alpha)$.
Lf	the Laplacian operator applied to the Boolean function f , defined by $Lf = \sum_{i=1}^n L_i f$ (or, the Ornstein–Uhlenbeck operator if f is a function on Gaussian space)
L_i	the i th coordinate Laplacian operator: $L_i f = f - \mathbf{E}_i f$
$\ln x$	$\log_e x$
$\log x$	$\log_2 x$
Maj_n	the majority function on n bits
$\mathbf{MaxInf}[f]$	$\max_i \{\mathbf{Inf}_i[f]\}$
$\mathbf{MaxInf}[f]$	$\max_i \{\mathbf{Inf}_i[f]\}$
$[n]$	$\{1, 2, 3, \dots, n\}$
\mathbb{N}	$\{0, 1, 2, 3, \dots\}$
$\mathbf{N}(0, 1)$	the standard Gaussian distribution
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$\mathbf{N}(0, 1)^d$	the distribution of d independent standard Gaussians; i.e., $\mathbf{N}(0, I_{d \times d})$
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$\mathbf{N}(\mu, \Sigma)$	for $\mu \in \mathbf{R}^d$ and $\Sigma \in \mathbf{R}^{d \times d}$ positive semidefinite, the d -variate Gaussian distribution with mean μ and covariance matrix Σ
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\mathbf{N}^+	$\{1, 2, 3, \dots\}$
$\mathbf{N}_{< m}$	$\{0, 1, \dots, m-1\}$
$N_\rho(x)$	when $x \in \{-1, 1\}^n$, denotes the probability distribution generating a string ρ -correlated to x
$N_\rho(z)$	when $z \in \mathbf{R}^n$, denotes the probability distribution of $\rho z + \sqrt{1-\rho^2} \mathbf{g}$ where $\mathbf{g} \sim \mathbf{N}(0, 1)^n$
$\mathbf{NS}_\delta[f]$	the noise sensitivity of f at δ ; i.e., $\frac{1}{2} - \frac{1}{2} \mathbf{Stab}_{1-2\delta}[f]$
\mathbf{OR}_n	the logical OR function on n bits: True unless all inputs are False
Φ	the standard Gaussian cdf, $\Phi(t) = \int_{-\infty}^t \phi(z) dz$
ϕ	the standard Gaussian pdf, $\phi(z) = \frac{1}{\sqrt{2\pi}} e^{-z^2/2}$
φ_A	the density function for the uniform probability distribution on A ; i.e., $1_A / \mathbf{E}[1_A]$
ϕ_α	given functions $\phi_0, \dots, \phi_{m-1}$ and a multi-index α , denotes $\prod_{i=1}^n \phi_{\alpha_i}$
$\bar{\Phi}$	the standard Gaussian complementary cdf, $\bar{\Phi}(t) = \int_t^\infty \phi(z) dz$
$\pi^{\otimes n}$	if π is a probability distribution on Ω , denotes the associated product probability distribution on Ω^n
$\pi_{1/2}$	the uniform distribution on $\{-1, 1\}$
π_p	the “ p -biased” distribution on bits: $\pi_p(-1) = p$, $\pi_p(1) = 1-p$
$\mathbf{Pr}_{\pi_p}[\cdot]$	an abbreviation for $\mathbf{Pr}_{x \sim \pi_p}[\cdot]$
\mathbf{R}	the real numbers
$\mathbf{R}^{\geq 0}$	the nonnegative real numbers
$\mathbf{RDT}(f)$	zero-error randomized decision tree complexity of f
$\mathbf{RS}_A(\delta)$	the rotation sensitivity of A at δ ; i.e., $\mathbf{Pr}[1_A(\mathbf{z}) \neq 1_A(\mathbf{z}')]]$ for a $\cos \delta$ -correlated pair $(\mathbf{z}, \mathbf{z}')$

$\text{sens}_f(x)$	the number of pivotal coordinates for f at x
$\text{sgn}(t)$	+1 if $t \geq 0$, -1 if $t < 0$
S_n	the symmetric group on $[n]$
$\text{sparsity}(f)$	$\Pr_{\mathbf{x}}[f(\mathbf{x}) \neq 0]$
$\text{sparsity}(\hat{f})$	$ \text{supp}(\hat{f}) $
$\mathbf{Stab}_\rho[f]$	the noise stability of f at ρ : $\mathbf{E}[f(\mathbf{x})f(\mathbf{y})]$ where \mathbf{x}, \mathbf{y} are a ρ -correlated pair
$\text{supp}(\alpha)$	if α is a multi-index, denotes $\{i : \alpha_i \neq 0\}$
$\text{supp}(f)$	if f is a function, denotes the set of inputs where f is nonzero
T_ρ	the noise operator: $T_\rho f(x) = \mathbf{E}_{\mathbf{y} \sim N_\rho(x)}[f(\mathbf{y})]$
T_ρ^i	the operator defined by $T_\rho^i f(x) = \rho f + (1 - \rho)\mathbf{E}_i f$
T_r	for $r \in \mathbf{R}^n$, denotes the operator defined by $T_{r_1}^1 T_{r_2}^2 \cdots T_{r_n}^n$
\mathcal{U}	the Gaussian isoperimetric function, $\mathcal{U} = \phi \circ \Phi^{-1}$
U_ρ	the Gaussian noise operator: $U_\rho f(z) = \mathbf{E}_{\mathbf{z}' \sim N_\rho(z)}[f(\mathbf{z}')]]$
$\mathbf{Var}[f]$	the variance of f , $\mathbf{Var}[f] = \mathbf{E}[f^2] - \mathbf{E}[f]^2$
\mathbf{Var}_i	defined by $\mathbf{Var}_i f(x) = \mathbf{Var}_{\mathbf{x}_i}[f(x_1, \dots, x_{i-1}, \mathbf{x}_i, x_{i+1}, \dots, x_n)]$
$\text{vol}_\gamma(A)$	$\Pr_{\mathbf{z} \sim N(0,1)^n}[\mathbf{z} \in A]$, the Gaussian volume of A
$\mathbf{W}^k[f]$	the Fourier weight of f at degree k
$\mathbf{W}^{>k}[f]$	the Fourier weight of f at degrees above k
$\mathbf{x} \sim A$	the random variable \mathbf{x} is chosen uniformly from the set A
$x^{(i \rightarrow b)}$	the string $(x_1, \dots, x_{i-1}, b, x_{i+1}, \dots, x_n)$
$x^{\oplus i}$	$(x_1, \dots, x_{i-1}, -x_i, x_{i+1}, \dots, x_n)$
$\mathbf{x} \sim \varphi$	the random variable \mathbf{x} is chosen from the probability distribution with density φ
x^S	$\prod_{i \in S} x_i$, with the convention $x^\emptyset = 1$
$\mathbf{x} \sim \{-1, 1\}^n$	the random variable \mathbf{x} is chosen uniformly from $\{-1, 1\}^n$
(y, z)	if $J \subseteq [n]$, $y \in \{-1, 1\}^J$, $z \in \{-1, 1\}^{\bar{J}}$, denotes the natural composite string in $\{-1, 1\}^n$
\mathbf{Z}	the additive group of integers modulo m
$\overline{\mathbf{Z}}_m^n$	the group indexing the Fourier characters of functions $f : \mathbf{Z}_m^n \rightarrow \mathbb{C}$