

---

Please print single-sided with each problem on its own pages. List any collaborators or sources (including yourself) at the end of your submission.

## 0 (optional ungraded practice)

(a) Prove the following languages are not regular:

- i.  $\{a+b=c \mid a, b, c \in \{0, 1\}^* \text{ and interpreted as binary representations of integers, } a + b = c\}$
- ii.  $\{1^p \mid p \text{ prime}\}$

(b) Give regular expressions for the following languages:

- i.  $\{x \in \{c, a, b\}^* \mid x \text{ contains the substring "cab"}\}$
- ii.  $\{x \in \{a, b, c\}^* \mid x \text{ contains an even number of } b\text{'s}\}$
- iii.  $\{x \in \{a, b, c\}^* \mid \text{every } a \text{ in } x \text{ is followed by a } b, \text{ and } x \text{ ends with a } c\}$

## 1 Deterministic Infinite Automata

We define a DIA =  $(Q, \Sigma, q_0, \delta, F)$  identically as for DFAs, except the set  $Q$  of states must be **infinite**. What set of languages do DIAs recognize? Prove your claim.

## 2 Blow-Up

Using the Rabin-Scott powerset construction from lecture, an NFA with  $n$  states can be made into an equivalent DFA with  $2^n$  states, but this is usually much larger than the minimal equivalent DFA.

Prove that for all  $n > 0$ , there is an  $n$ -state NFA whose minimal equivalent DFA has at least  $2^{n-1}$  states.

## 3 Converse Pumping Lemma

Consider the language  $A = \{a^i b^j c^k \mid i, j, k \geq 0 \text{ and if } i = 1 \text{ then } j = k.\}$

(a) Show that  $A$  is not regular.

(b) Show that the pumping lemma is insufficient to prove  $A$  is not regular, that is, we have some  $P$  where if  $w \in A$  and  $|w| \geq P$ , we can write  $w = xyz$  with  $|y| > 0$ ,  $|xy| \leq P$ , and for every  $i \geq 0$ ,  $xy^i z \in A$ .

(c) Do (a) and (b) contradict the pumping lemma? Explain.