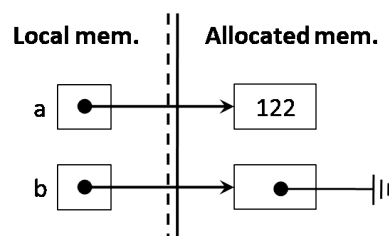


## Pointer manipulations

Pointer manipulations are counterintuitive at first, but with just a little bit of practice, they will become second nature. A key insight is the following:

*When you set a pointer equal to another pointer, you make the first pointer point to where the second pointer points.*

Consider the following memory diagram, where `a` has type `int*` and `b` has type `int**`:



Draw the memory diagram after executing the code line

```
*b = a;
```

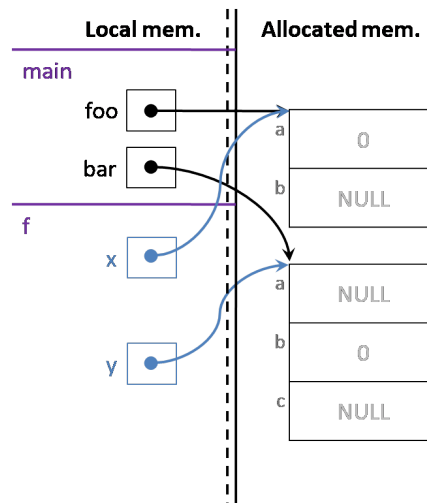
## A wild struct appears

Suppose we have the following in a file:

```

1 struct X {
2     int    a;
3     struct Y* b;
4 };
5
6 struct Y {
7     int*    a;
8     int     b;
9     struct X* c;
10 };
11
12 void f(struct X* x, struct Y* y) {
13     x->b = y;
14     y->c = x;
15     y->c->a = 15
16     int** d = alloc(int*);
17     *d = alloc(int);
18     x->b->a = *d;
19     *(y->a) = x->a * 8 + 2;
20     x->b->b = 1000 * x->a + **d;
21     x = NULL;
22     y->c = NULL;
23     return;
24 }
25
26 int main() {
27     struct X* foo = alloc(struct X);
28     struct Y* bar = alloc(struct Y);
29
30     return 0;
31 }

```



## Checkpoint 0

Fill out the state of the memory. What's the value of `bar->b`? (For your own sanity, draw a picture!)

## Stack and queue interfaces

Here's the **stack interface** discussed in lecture. It exposes the type `stack_t` and four functions:

```
// typedef _____* stack_t;    /* Abstract type of stacks          */

bool stack_empty(stack_t S)      /* Check if stack S is empty,      0(1) */
/*@requires S != NULL; @*/ ;

stack_t stack_new()             /* Create a new empty stack,       0(1) */
/*@ensures \result != NULL; @*/
/*@ensures stack_empty(\result); @*/ ;

void push(stack_t S, string x) /* Add item x at the top of stack S, 0(1) */
/*@requires S != NULL; @*/
/*@ensures !stack_empty(S); @*/ ;

string pop(stack_t S)          /* Remove and return the top of stack S, 0(1) */
/*@requires S != NULL; @*/
/*@requires !stack_empty(S); @*/ ;
```

The **queue interface** exposes the type `queue_t` and four similar functions:

```
// typedef _____* queue_t;    /* Abstract type of queues          */

bool queue_empty(queue_t Q)     /* Check if queue Q is empty,      0(1) */
/*@requires Q != NULL; @*/ ;

queue_t queue_new()            /* Create a new empty queue,       0(1) */
/*@ensures \result != NULL; @*/
/*@ensures queue_empty(\result); @*/ ;

void enq(queue_t Q, string e) /* Add item e at the back of queue Q, 0(1) */
/*@requires Q != NULL; @*/
/*@ensures !queue_empty(Q); @*/ ;

string deq(queue_t Q)         /* Remove and return the front of queue Q, 0(1) */
/*@requires Q != NULL; @*/
/*@requires !queue_empty(Q); @*/ ;
```

## Checkpoint 1

Write a function to reverse a queue using only functions from the stack and queue interfaces.

```

1 void reverse(queue_t Q)
2 //@requires _____;
3 {
4 _____ // create temp data structure
5 while ( _____ ) {
6 _____
7 }
8 while ( _____ ) {
9 _____
10 }
11 }

```

## Checkpoint 2

Write a *recursive* function to count the size of a stack. You may not destroy the stack in the process — the stack’s elements (and order) must be the same before and after calling this function. Assume the stack contains elements of type **string**.

```

int size(stack_t S)
//@requires _____;
{
_____
_____
_____
_____
_____
}

```

## Checkpoint 3

Why couldn’t this function be used in contracts in C0? Hint: Contracts in C0 can’t have side effects.