

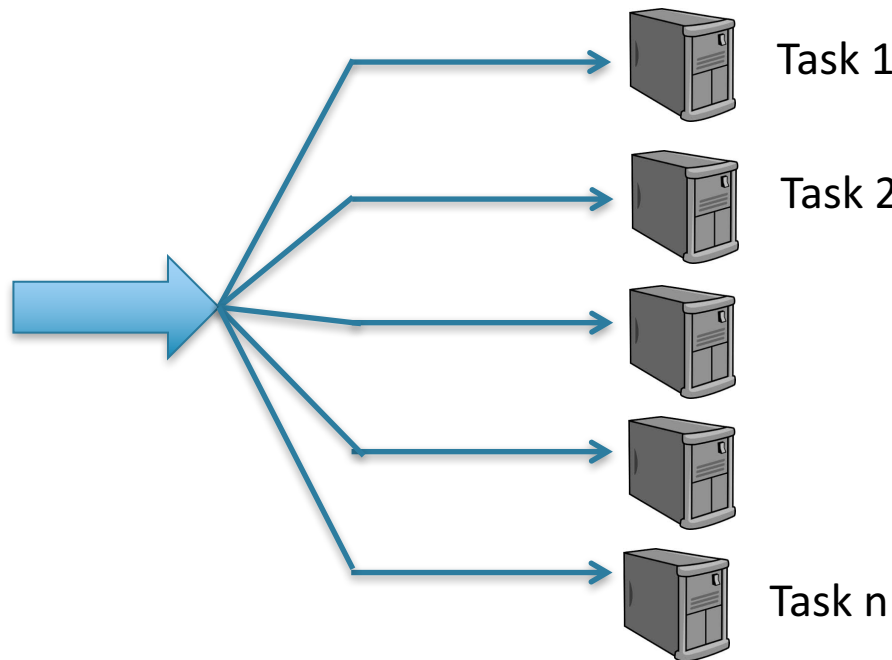
Using Straggler Replication to Reduce Latency in Large-scale Parallel Computing

Da Wang, Gauri Joshi, Gregory Wornell



Problem: Stragglers in Parallel Computing

- A job with hundreds of parallel tasks
- Machine response time can vary due to virtualization, congestion etc.
- The slowest tasks are the bottleneck in job completion

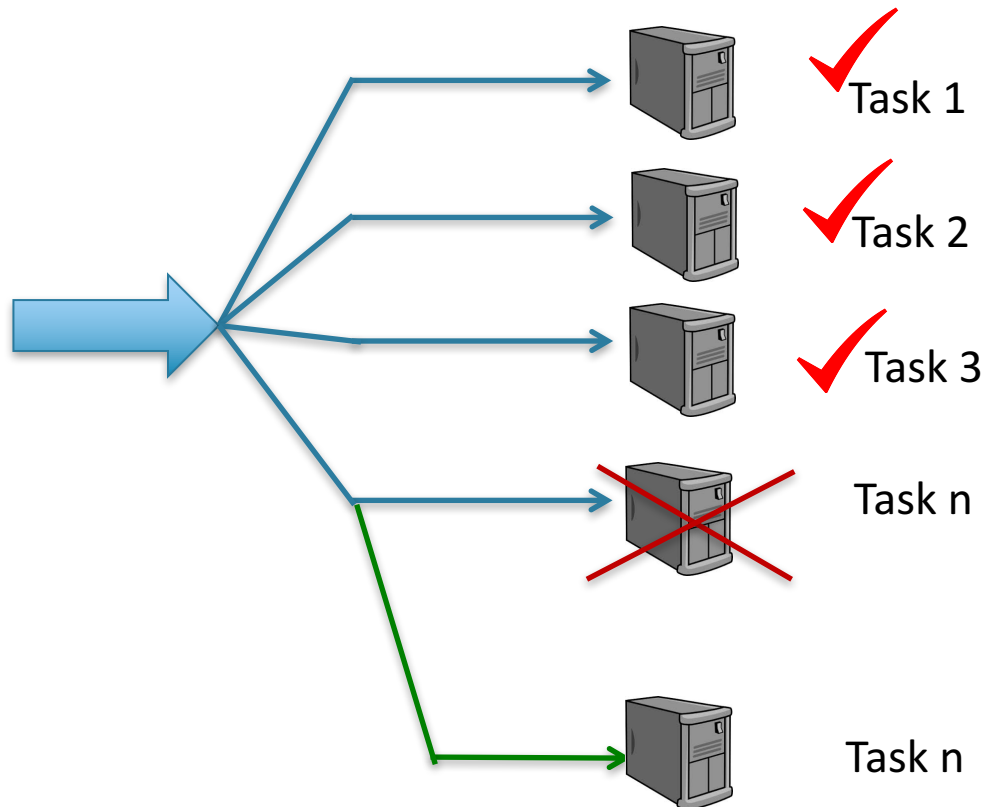


[Dean “Tail at Scale” 2013]

Latency	50%ile	99%ile
1 task finishes	1ms	10ms
All tasks finish	40ms	140ms

Solution: Replication of Stragglers

Re-run the stragglers when p fraction of tasks are remaining



Related Previous Work

Task Replication in Systems Literature

- First used in MapReduce [Dean 2008] via back-up tasks
- Further developed in [Zaharia 2008], [Ananthanarayanan 2010] etc

Used in practice, but little theoretical analysis so far

Our Contributions

- Provide design insights on how to schedule task replication to reduce delay, with efficient use of additional resources

[1] D. Wang, G. Joshi, G. Wornell, " **Efficient Task Replication for Fast Response Times in Parallel Computation** ", SIGMETRICS 2014

[2] D. Wang, "Computing with Unreliable Resources", PhD Thesis, MIT, 2014

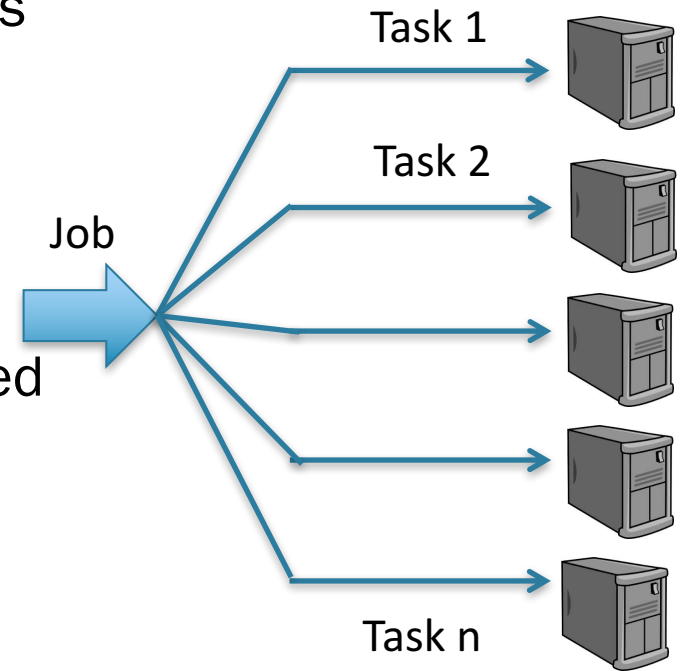
[3] G. Joshi, "Efficient Redundancy Techniques to Reduce Delay in Cloud Systems", PhD Thesis, MIT 2016

System Model

- A job with n parallel tasks, n is large
- Finish time of a task, $X \sim F_X$, i.i.d. across machines.

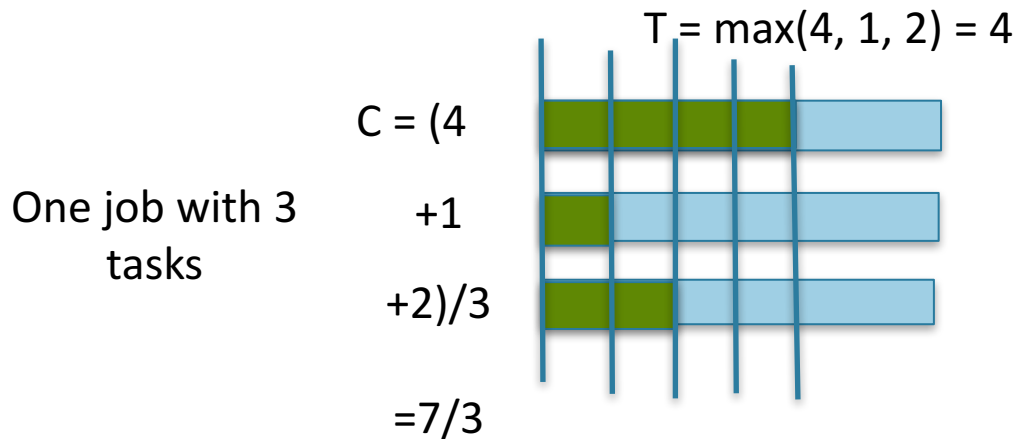
Remark on the i.i.d assumptions:

- From cloud user's point of view, all rented machines are approx. identical.



Performance Metrics

- Expected Latency $E[T]$ = Expected Time when all tasks finish
- Expected Cost $E[C]$ = Expected total machine time spent, normalized by # of tasks



Remark on cost metric

- There could be other costs – network, memory usage, etc

Outline

Analysis of $E[T]$ and $E[C]$ using Extreme Value Theory

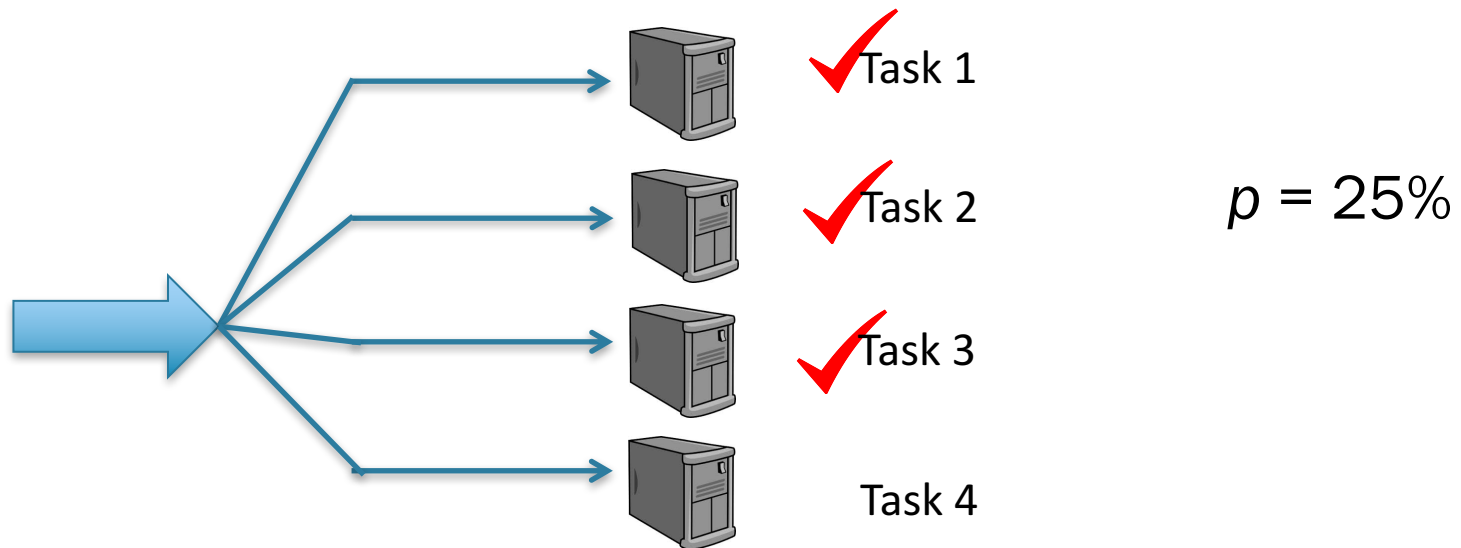
- Tail behavior of F_x is a key factor affecting the $E[T]$ - $E[C]$ trade-off

Heuristic Algorithm to find best replication strategy

- Compare with back-up tasks in MapReduce using Google trace data

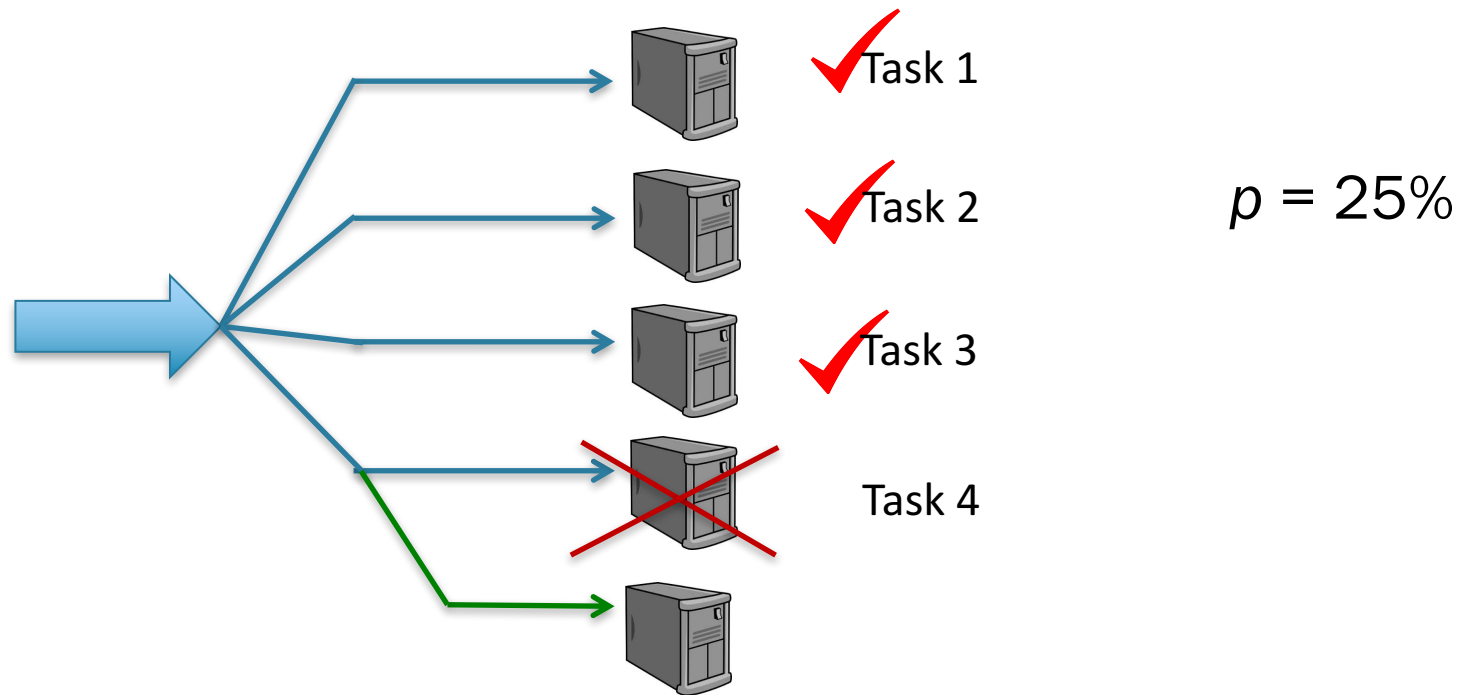
Replication Policy: When to re-run?

- Re-run the stragglers when p fraction of tasks are left



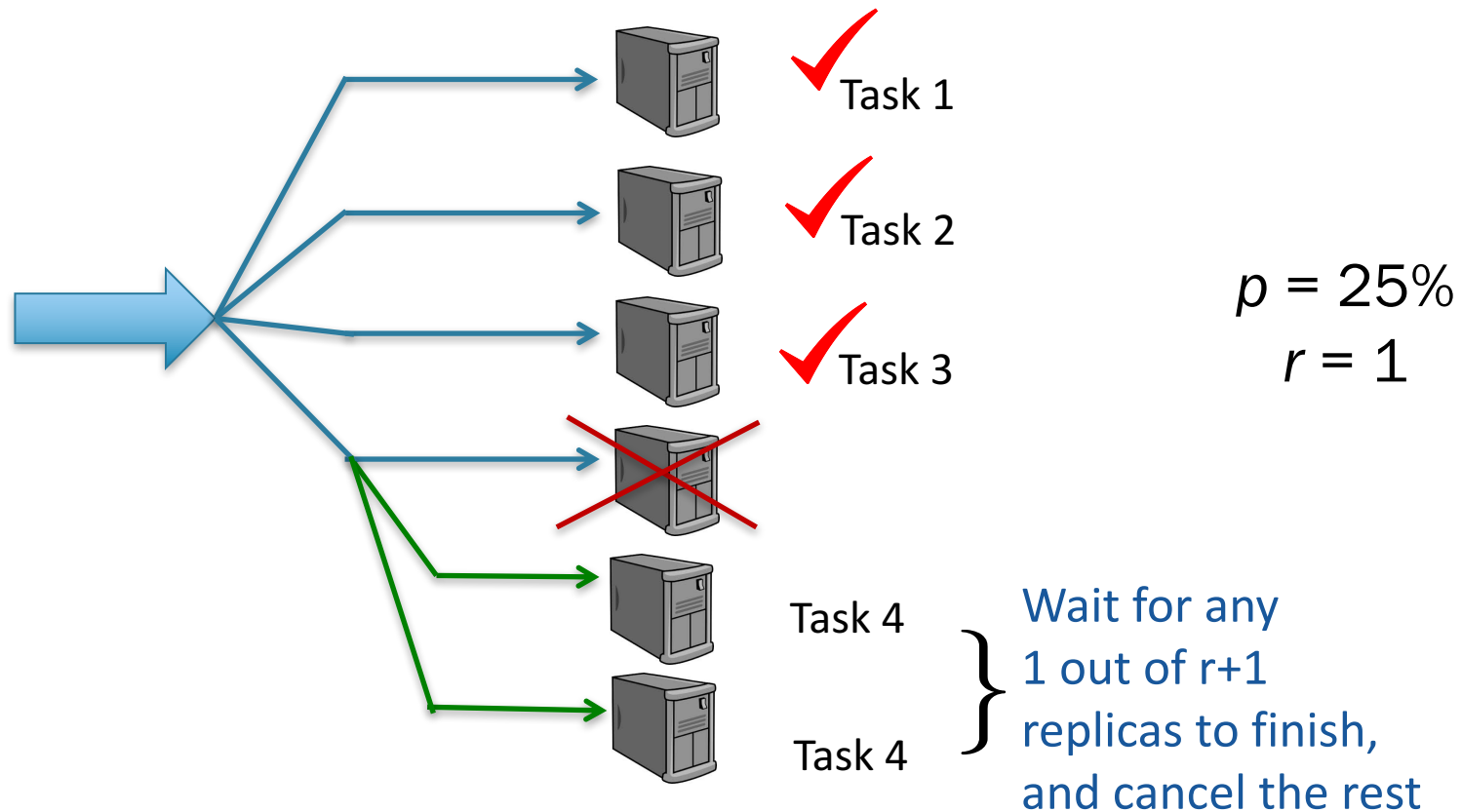
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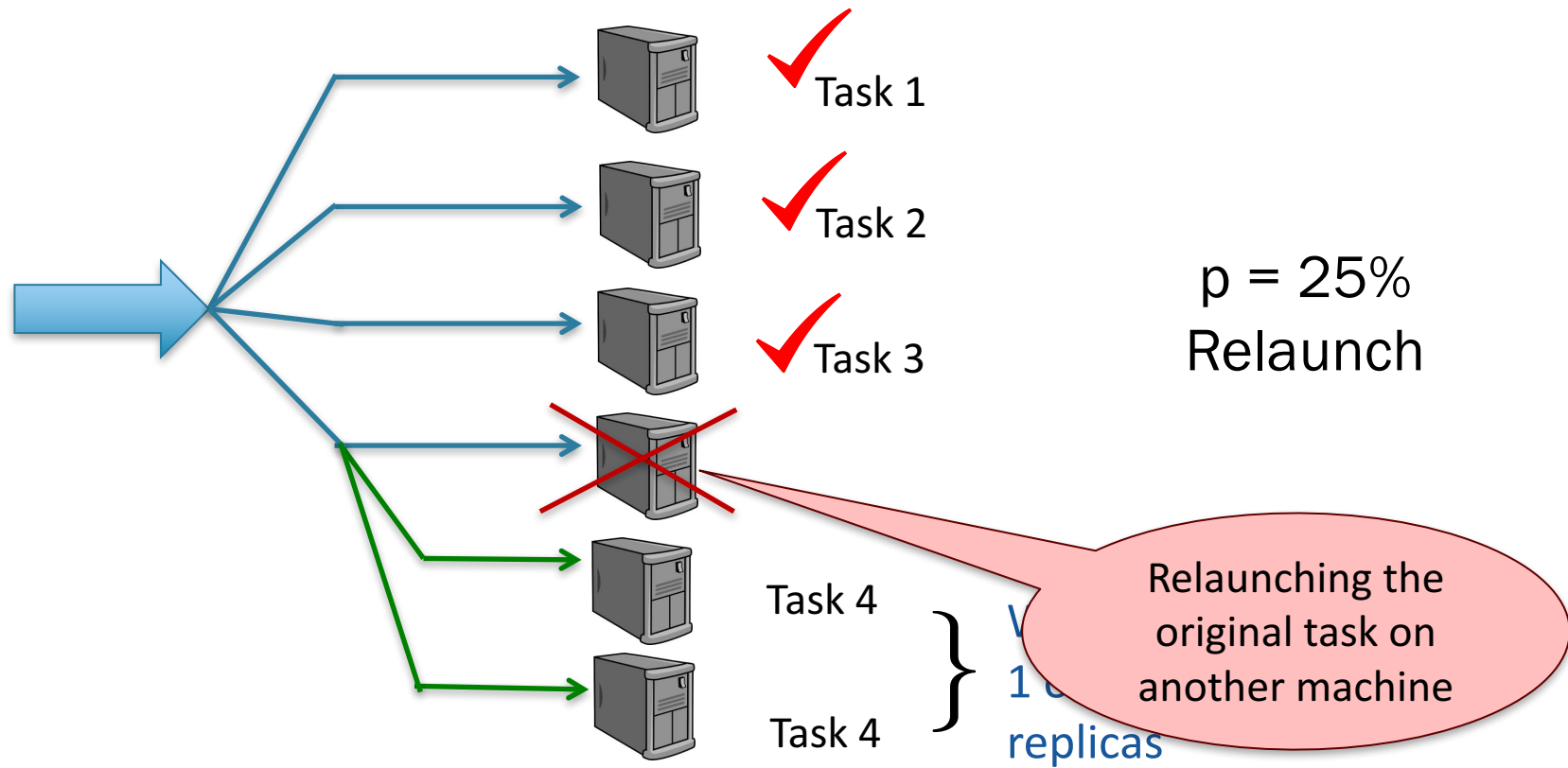
Replication Policy: How many replicas?

- Re-run the stragglers when p fraction of tasks are left
- Run r additional replicas



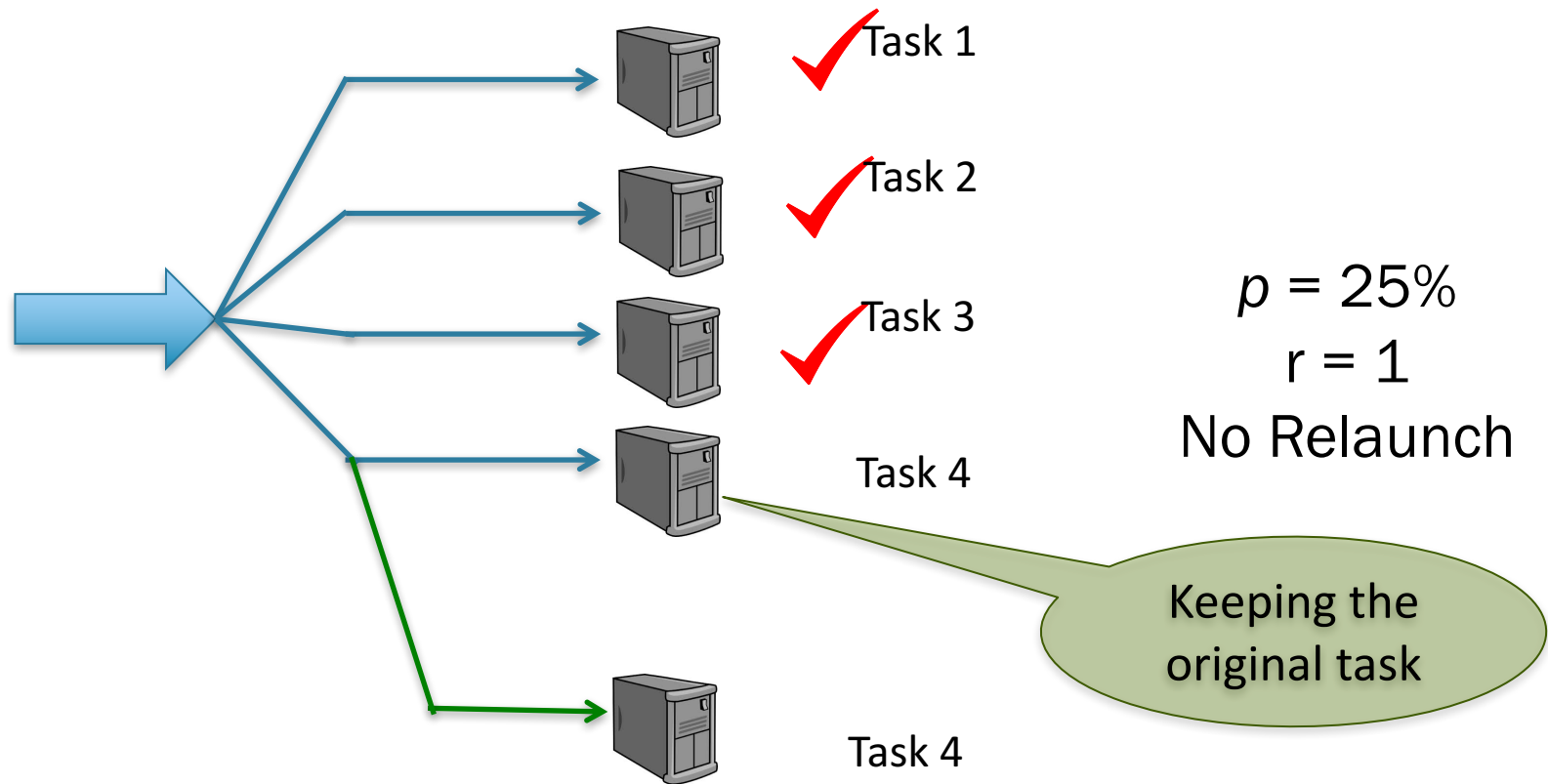
Replication Policy: Relaunch or not?

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Replication Policy: Relaunch or not?

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Problem Formulation

Given n tasks, and task finish time distribution F_X ,

Design Parameters

- p : Fraction of tasks left when we replicate
- r : Number of additional replicas
- Relaunch original straggling task or not

Performance Metrics

- Latency $E[T]$
- Cost $E[C]$

Evaluating Expected Latency $E[T]$

- Wait for $(1-p)n$ tasks to finish
- Launch replicas of the pn stragglers
 - Time for 1 out of $r+1$ copies to finish $Y \sim F_Y = g(F_X, r, \text{kill/keep})$
 - For e.g. $r=1$ with task-killing $\rightarrow (1-F_Y) = (1-F_X)^2$

$$T = X_{(1-p)n:n} + Y_{pn:pn}$$

Notation $X_{k:n}$:
 k^{th} smallest of n i.i.d.
rvs X_1, X_2, \dots, X_n

Wait for $(1-p)n$ out of n tasks to finish

Maximum of finish times of the pn stragglers after replication

Evaluating Expected Latency $E[T]$

$$\mathbb{E}[T] = \mathbb{E}[X_{(1-p)n:n}] + \mathbb{E}[Y_{pn:pn}]$$

Central Value Theorem $n \rightarrow \infty$

Extreme Value Theorem $n \rightarrow \infty$

$F_X^{-1}(1-p)$

Different behavior for Exponential, Light or Heavy tailed Y

Asymptotic approx for $n \rightarrow \infty$ is close to simulation even for $n \sim 300$

Exercise: Task Execution Time $X \sim \text{Exp}(\mu)$

$$\mathbb{E}[T] = \mathbb{E}[X_{(1-p)n:n}] + \mathbb{E}[Y_{pn:pn}]$$

Central Value Theorem $n \rightarrow \infty$

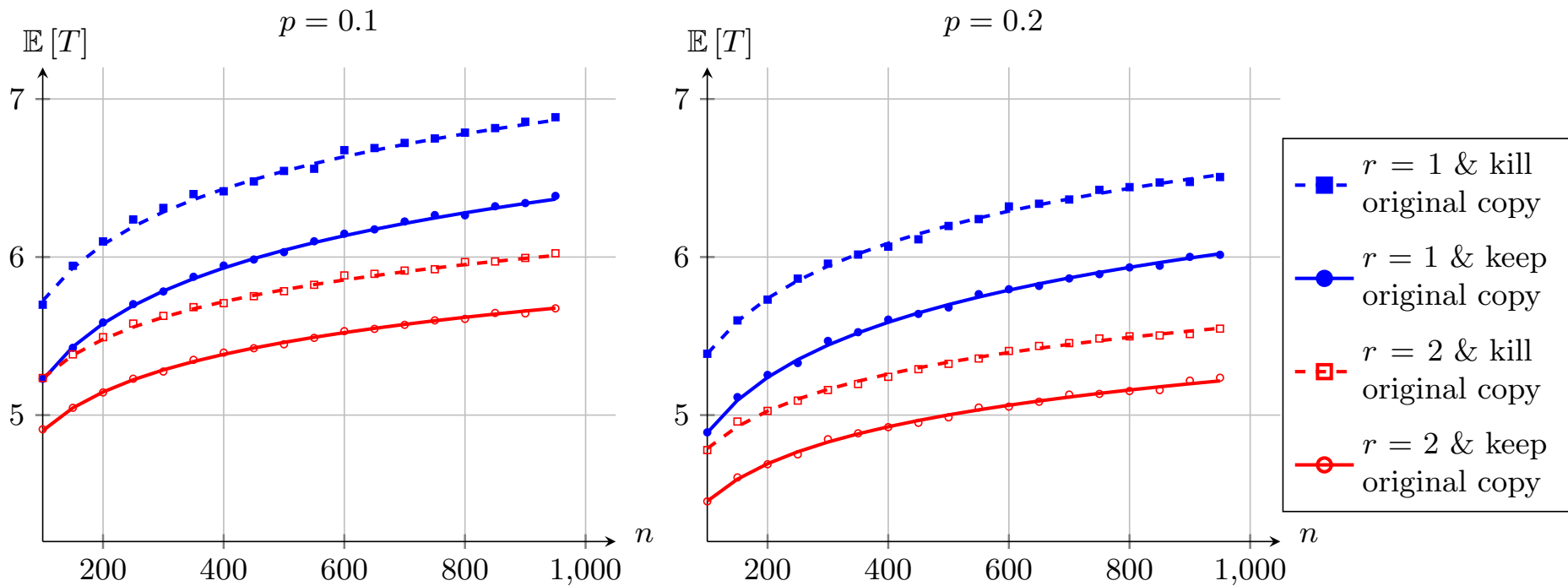
Extreme Value Theorem $n \rightarrow \infty$

$F_X^{-1}(1-p)$

Different behavior for Exponential, Light or Heavy tailed Y

Comparing Theoretical Analysis with Simulations

$X \sim 1 + \text{Exp}(1)$, and $n = 400$

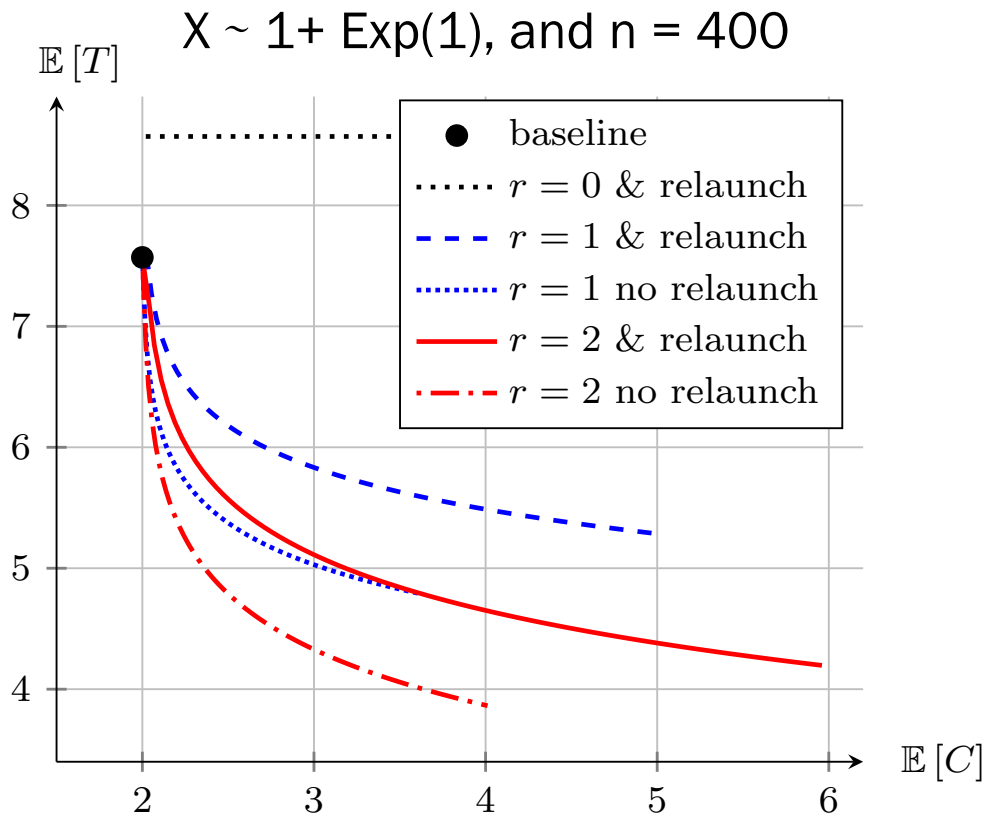


Evaluating Expected Cost $\mathbb{E}[C]$

$$\begin{aligned}\mathbb{E}[C] &= \frac{1}{n} \sum_{i=1}^{(1-p)n} \mathbb{E}[X_{i:n}] + \frac{np}{n} \mathbb{E}[T^{(1)}] + \frac{1}{n} \sum_{j=1}^{pn} (r+1) \mathbb{E}[Y] \\ &= \int_0^{1-p} F_X^{-1}(h) dh + pF_X^{-1}(1-p) + (r+1)p\mathbb{E}[Y] + O(1/n)\end{aligned}$$

By Central
Value Theorem

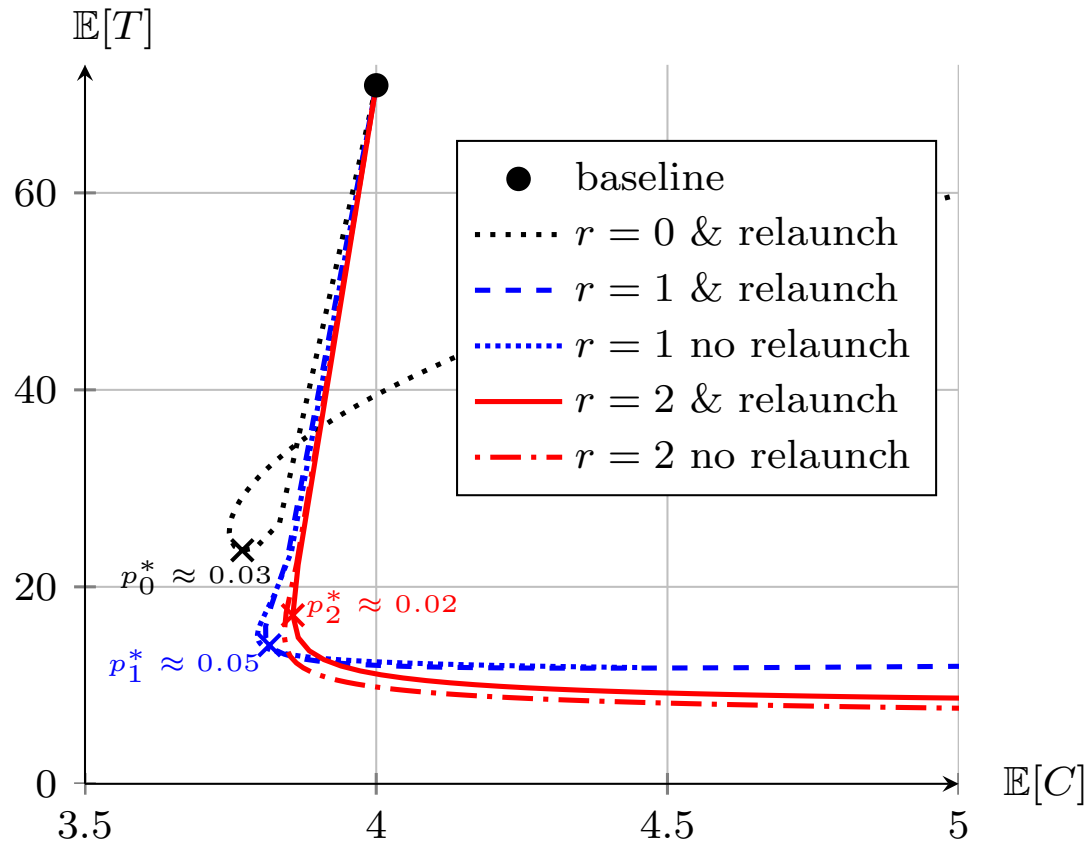
Case: Shifted Exponential (Exp. tail)



- Increasing p and r reduces latency but increases cost
- Killing a straggling task never helps!

Case: Pareto (Heavy tail)

$X \sim \text{Pareto}(2, 2)$ and $n = 400$



Latency and cost both reduce for small p !

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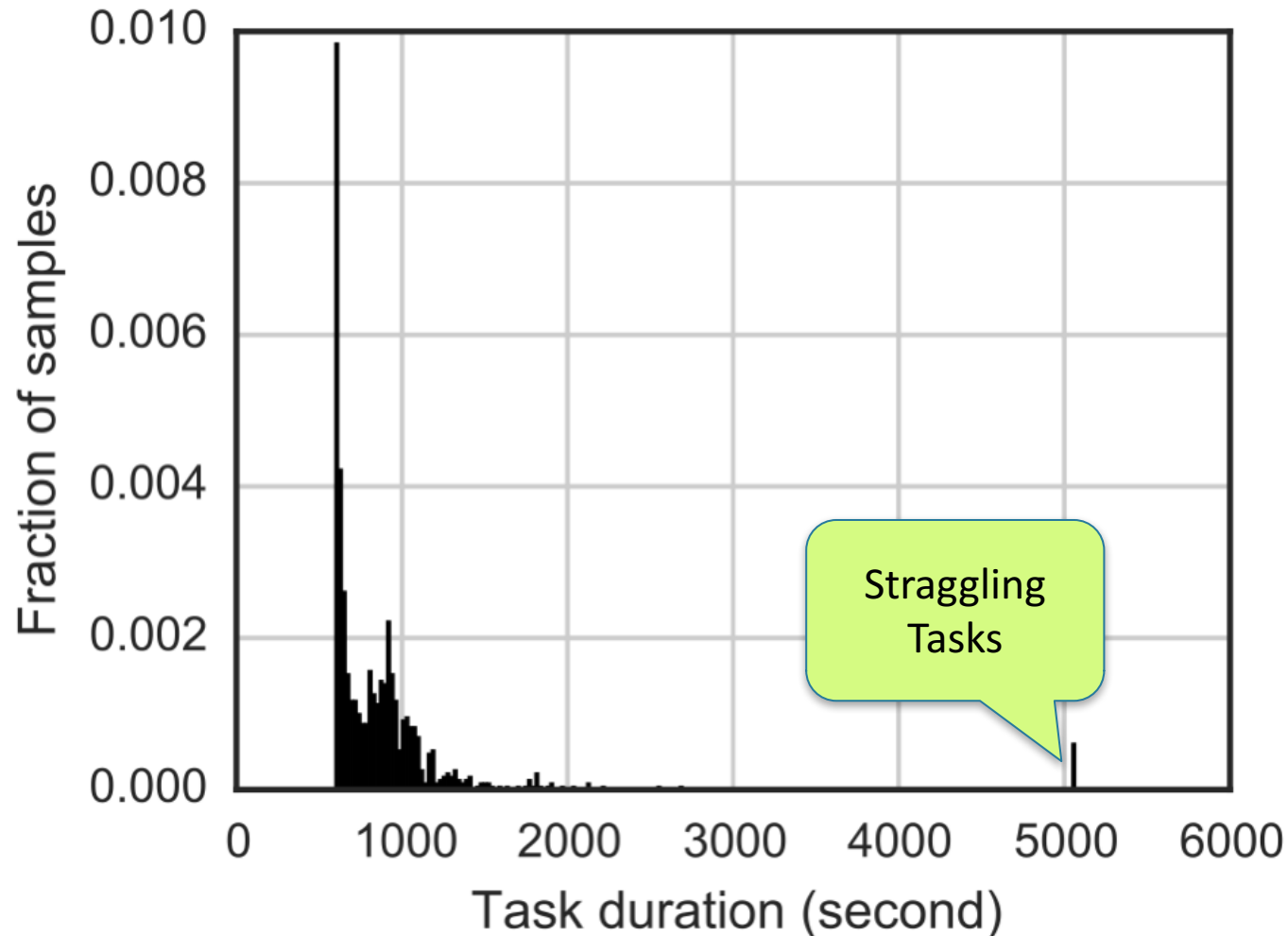
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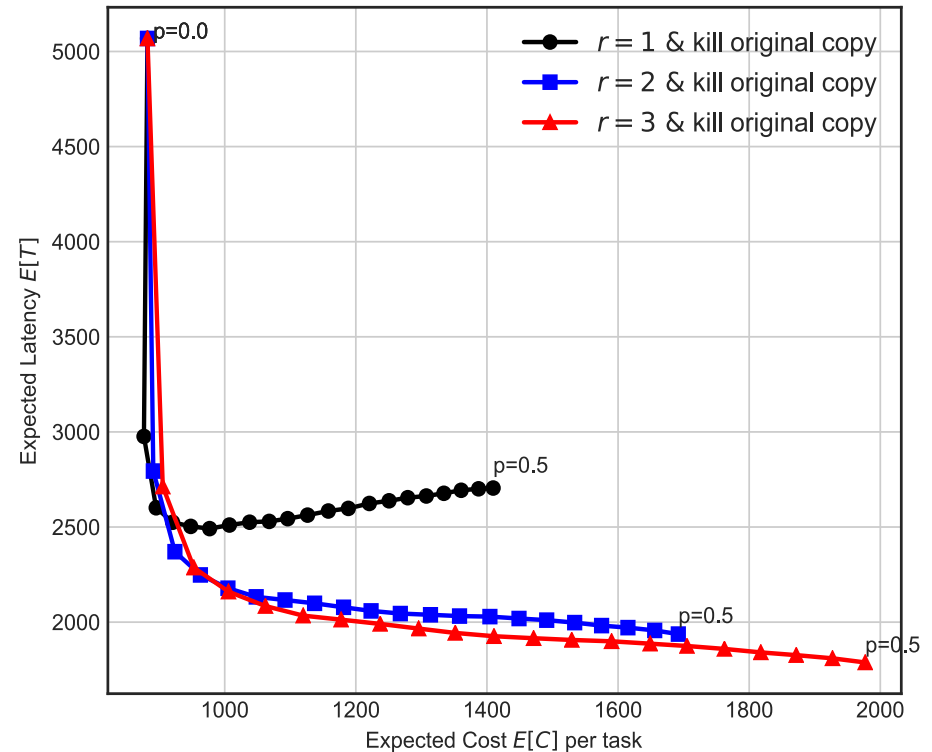
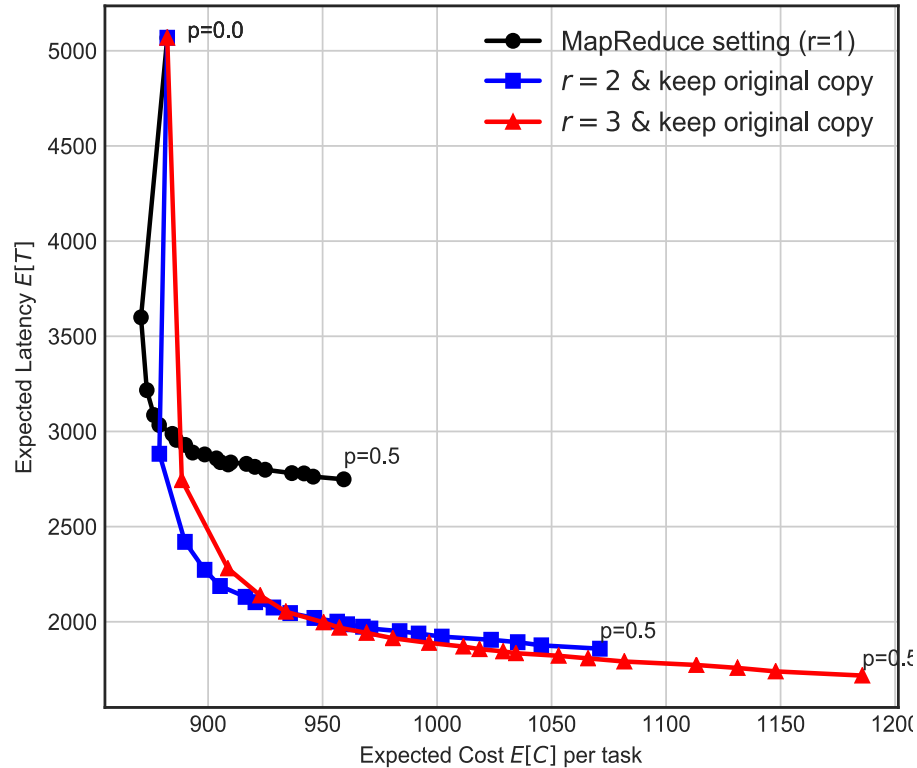
Example: Job with 1026 tasks

Google Cluster Data



Simulations using Google Cluster Data

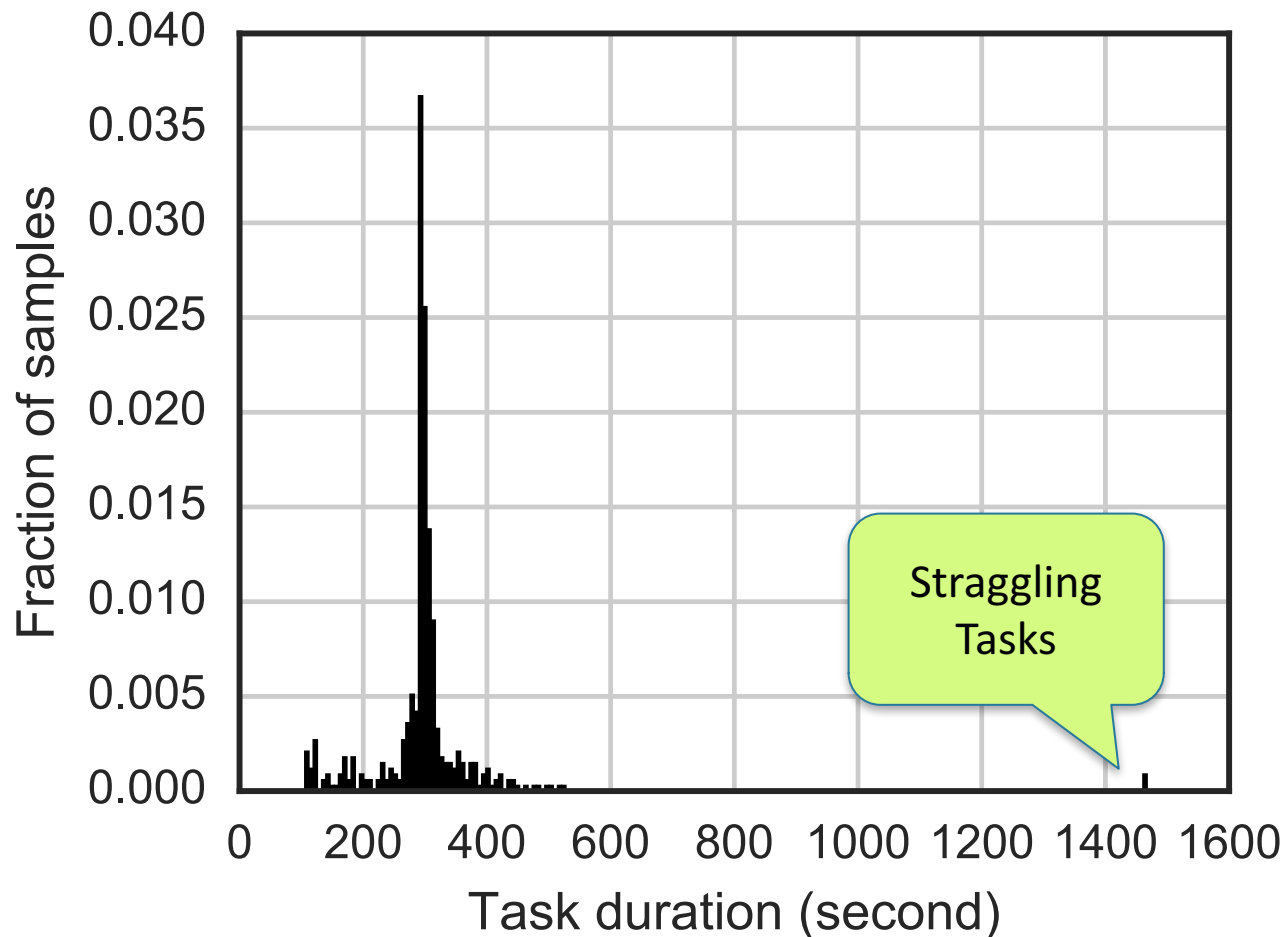
Latency-Cost Trade-off



Careful choice of replication strategy can be better than the default in MapReduce

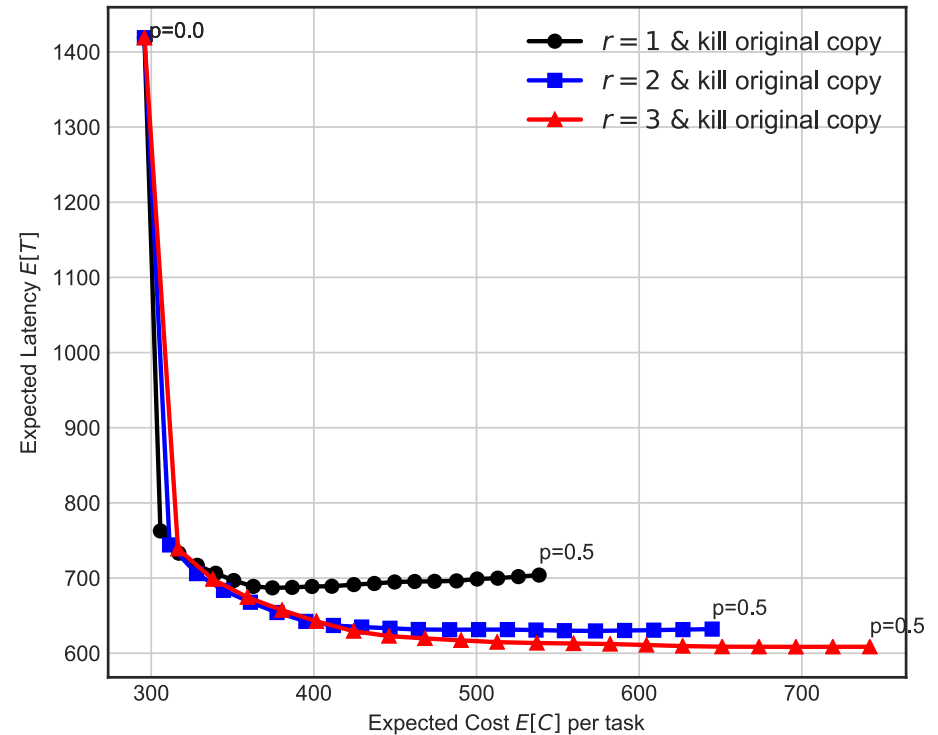
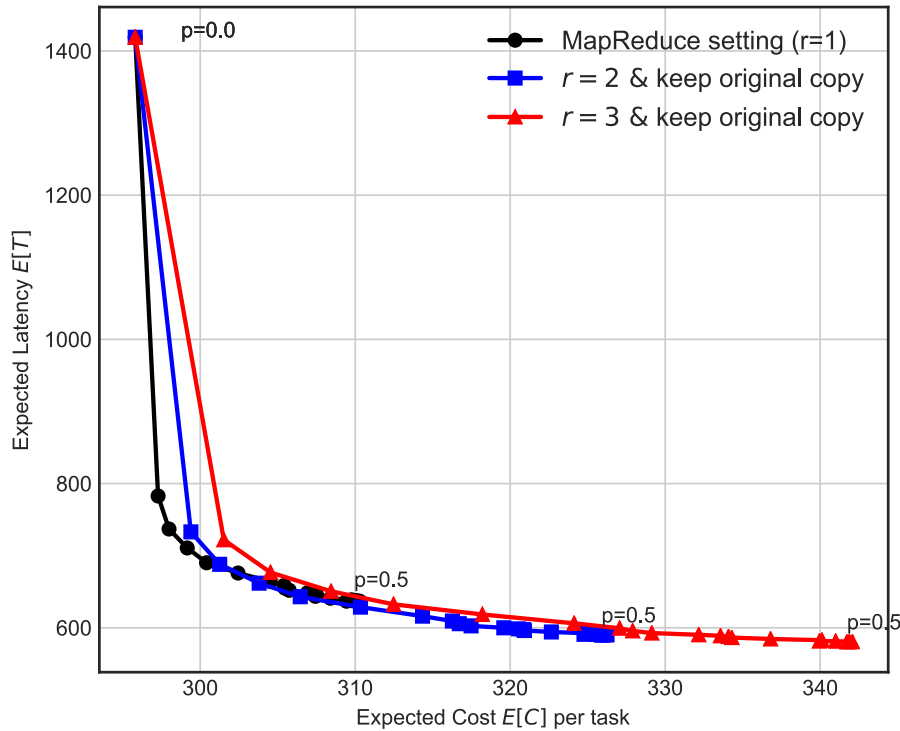
Example: Job with 488 tasks

Google Cluster Data



Simulations using Google Cluster Data

Latency-Cost Trade-off



Heuristic Search of the Best Strategy

May be hard to use our analysis to optimize the strategy for any F_x

- Analysis of $E[T]$ and $E[C]$ after replication can be hard

ESTIMATION

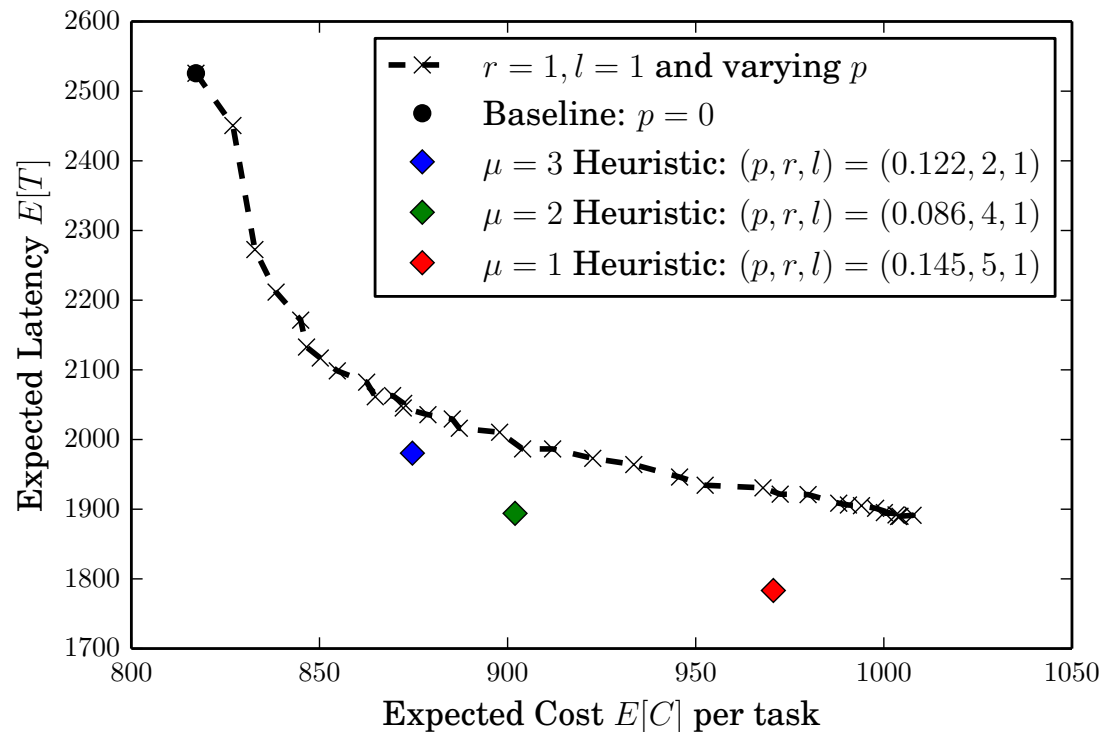
- Estimate F_x from traces of task execution time
- Use empirical F_x to estimate $J = E[T] + \mu E[C]$ for given p , r , relaunch/not

HEURISTIC ALGORITHM

1. For given p , choose r , and the kill/keep strategy that minimizes J
2. Perform gradient descent on p

Heuristic Algo: Resulting $E[T]$ and $E[C]$

- Run heuristic algorithm with different μ , to minimize $J = E[T] + \mu E[C]$
- $r = 1$, without relaunch ($l=1$): Back-up tasks option in MapReduce



Concluding Remarks

SUMMARY

- Tail behavior is important in choosing the right policy
 - Pareto, Shifted Exponential etc.
- Heuristic algorithm to find good replication policy given traces of execution time

RELATED AND FUTURE WORK

- Queueing of jobs (next class)
- Online algorithm to learn F_x and schedule simultaneously